

Preserve System Integrity for Your Business With IBM Replication Solutions for Business Continuity

Part 2 of 2 Session # 9948

Jeff Suarez IBM Corporation jrsuarez@us.ibm.com Glenn Wilcock IBM Corporation wilcock@us.ibm.com

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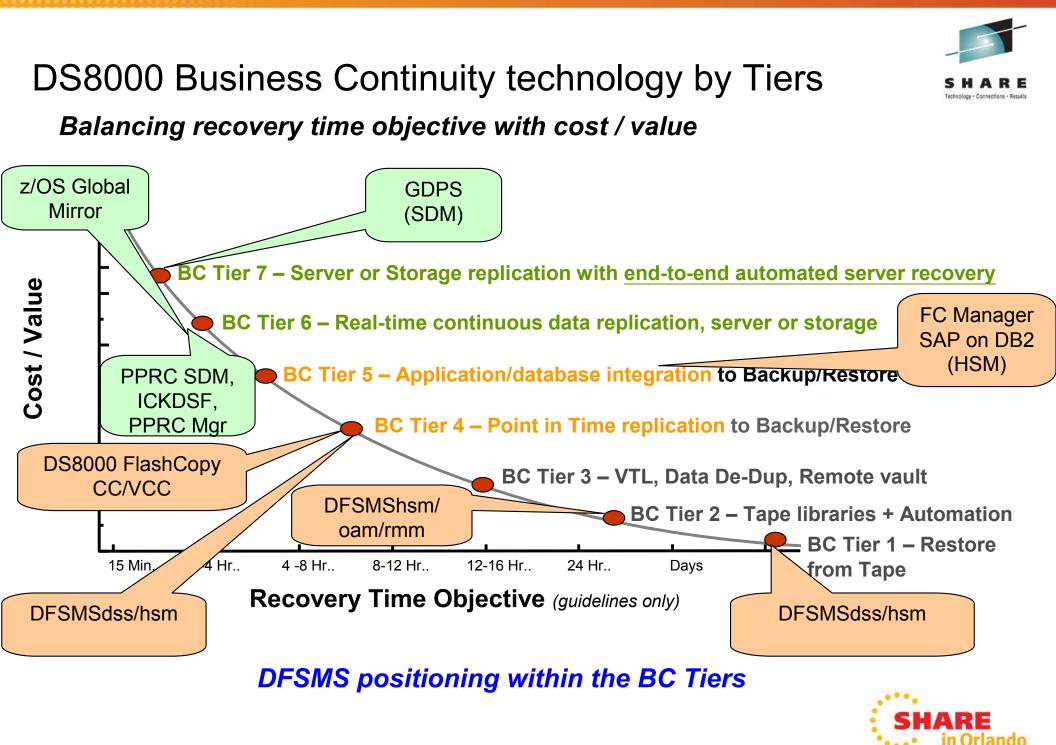
Agenda

- Business Continuity Overview
- DFSMSdss
 - Overview
 - Commands
 - FlashCopy

Introduction to Continuous Data Protection

- Overview
- Types
- DB2 Solution
- SMS Copy Pool
- DFSMShsm
 - Fast Replication Backup
 - Fast Replication Recover
- DB2 Overview







DFSMSdss Terms

- Physical Processing
 - Volume view for task
- Logical Processing
 - Dataset view for task
- Data Set Filtering
 - INCLUDE list consists of what data sets you want processed
 - Specific data set names (fully qualified), wildcards (*, **, %)
 - EXCLUDE list contains data sets you don't want processed
 - Specific data set names (fully qualified), wildcards (*, **, %)
 - BY Filtering
 - Check data set type, size, catalog status, SMS class, etc.
 - FILTERDD
 - DD Statement points to data set containing INCLUDE/EXCLUDE/BY filtering



DFSMSdss Interfaces



Batch JCL

• PGM=ADRDSSU

Application Programming Interface (API)

- DFSMSdss may be invoked by other programs for any of its functions (except for Stand Alone Restore)
- Invoking program can pass address of a User Interaction Module (UIM)
- DFSMdss calls the UIM at various exit points throughout processing
- UIM may direct DFSMSdss processing via these exit points

Cross Memory API

- Method to get DFSMSdss functionality without paying the penalty of having the DFSMSdss processing (and memory usage) in your own address space
- Client / Server model with a small client in the invoking program's address space, and a server in its own address space that attaches ADRDSSU tasks to fulfill the client's request
- May be invoked via Batch JCL (PGM=ADRXMAIA)



DFSMSdss Commands



- BUILDSA
- CGCREATED *
- COMPRESS
- CONSOLIDATE *
- CONVERTV
- COPYDUMP
- * Exploits Fast Replication function

- DEFRAG *
- PRINT
- RELEASE
- COPY *
- DUMP *
- RESTORE



SHARE Technology - Connections - Results

DFSMSdss FlashCopy

- Functions that exploit FlashCopy
 - COPY
 - Full volume
 - DUMPCONDITIONING available to provide online volume clone to be used for backup purposes
 - Data Set level
 - Logical or Physical
 - Physical allows copying back a set of files using FC on a volume basis
 - DataSet DUMPCONDITIONING
 - Can Rename data sets to a temporary name at COPY time, and at DUMP time have the ability to rename them back to its original name
 - DEFRAG and CONSOLIDATE
 - Default is to use FlashCopy if possible, can revert to using CC,VCC, and traditional I/O
 - Updates catalog and performs serializations
 - Volume and data set level





DFSMSdss FlashCopy Functions

- DFSMSdss FlashCopy Functions, Commands, Keywords
 - For COPY
 - FCNOCOPY and FCNOCOPYTOCOPY
 - FCFREEZE and CGCREATE
 - FCINCREMENTAL, FCINCREMENTALLAST
 - FCINCRVERIFY(REVERSE|NOREVERSE) and FCWAIT
 - FCSETGTOK(FAILRELATION)
 - FCTOPPRCPRIMARY
 - Preserve Mirror Options
 - FCFASTREVERSERESTORE and FCFULLVOLUMERELATION (New)
 - For DUMP
 - FCWITHDRAW (includes INIT and Space Release for SE Volumes)
 - Debugging Keywords for Fast Replication
 - DEBUG(FRMSG(DTL|SUMM|NO))
 - Valid for COPY, DEFRAG, and CONSOLIDATE





- Limits background copy to tracks that have changed since the previous Incremental FlashCopy
 - First Incremental FlashCopy results in full background copy and initiates change recording
 - Subsequent Incremental FlashCopies either:
 - Refresh the target: only the changed tracks being copied

or

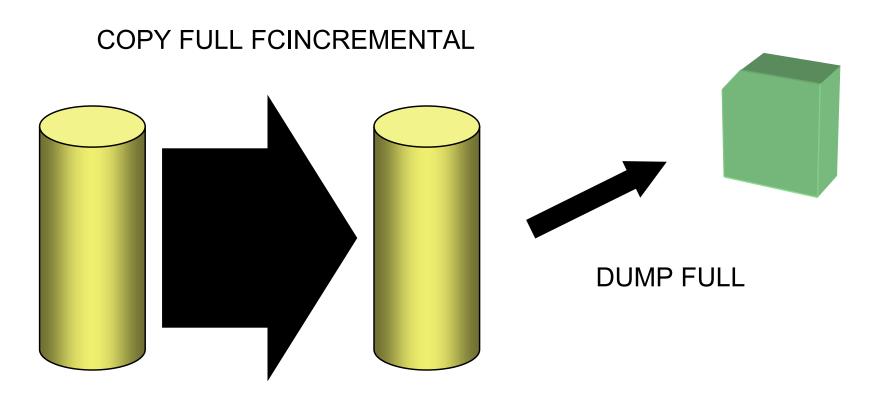
- Restore the source: changed tracks are copied back, removing updates
- Incremental relationships are persistent (remain after background copy is complete)
- Only one incremental relationship per volume
 - Can exist with other non-incremental relationships







Usage Scenario 1: Periodic Dumps to Tape



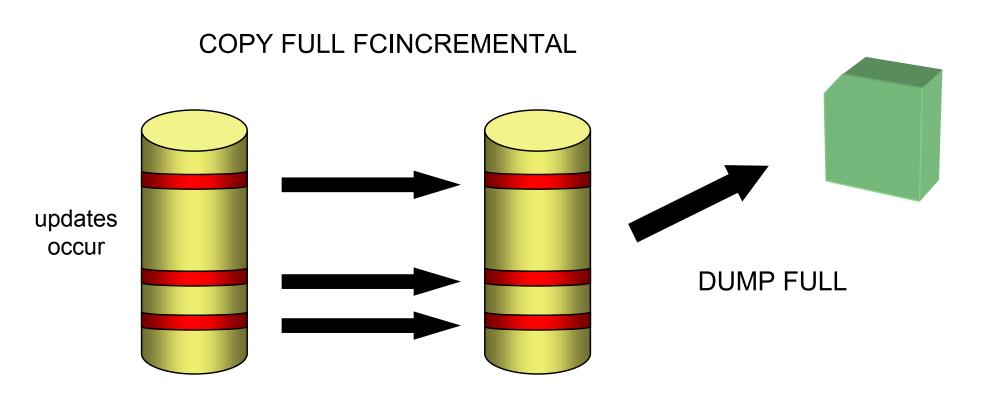
First time, all data is copied







Usage Scenario 1: Periodic Dumps to Tape

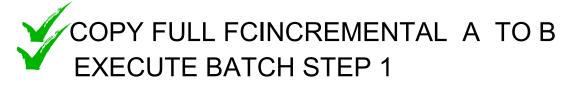


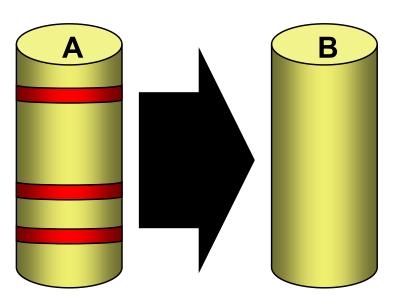
only updated tracks are copied





Usage Scenario 2: Batch Checkpoints



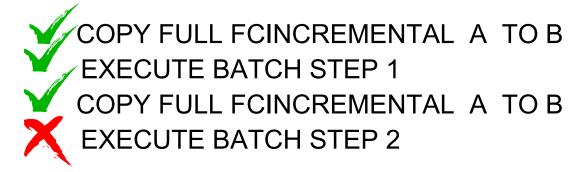


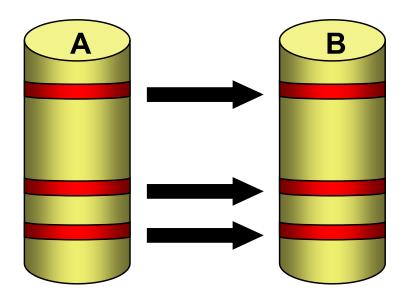
all data is copied





Usage Scenario 2: Batch Checkpoints



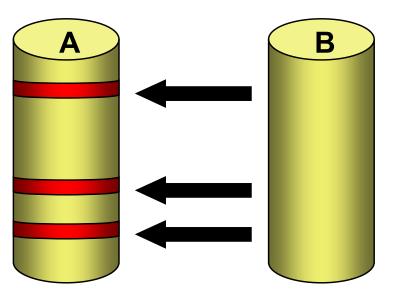


only updates are copied





Usage Scenario 2: Batch Checkpoints



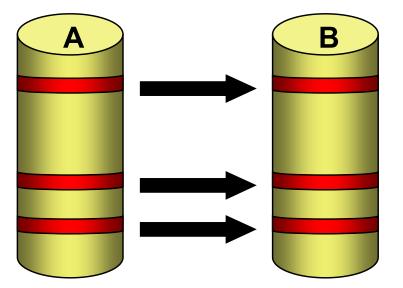
updates are removed

COPY FULL FCINCREMENTAL A TO B EXECUTE BATCH STEP 1 COPY FULL FCINCREMENTAL A TO B EXECUTE BATCH STEP 2 COPY FULL FCINCREMENTAL B TO A FCINCRVERIFY(REVERSE) FCWAIT(2) RESTART BATCH STEP 2





Usage Scenario 2: Batch Checkpoints



only updates are copied

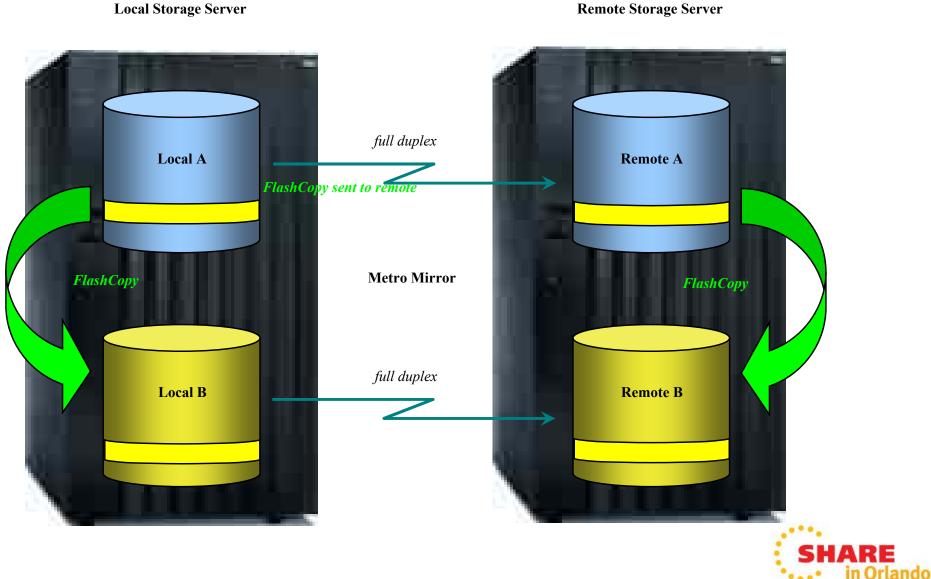
 COPY FULL FCINCREMENTAL A TO B EXECUTE BATCH STEP 1
 COPY FULL FCINCREMENTAL A TO B EXECUTE BATCH STEP 2
 COPY FULL FCINCREMENTAL B TO A FCINCRVERIFY(REVERSE) FCWAIT(2)
 RESTART BATCH STEP 2
 COPY FULL FCINCREMENTAL A TO B FCWAIT(2)
 EXECUTE BATCH STEP 3



FlashCopy to PPRC Primary using Preserve **Mirror**



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Remote Storage Server



Continuous Data Protection for DB2

zCDP for DB2





Business Continuity Overview

Business Continuity

Maintaining business operations in the event of an outage – with processes and infrastructure that are responsive, highly available and scalable

Three key characteristics

Recovery Time Objective (RTO)

• The *acceptable* amount of time you can afford to be without your data

Recovery Point Objective (RPO)

• The amount of data that can be acceptably recreated

✓ Backup Window Objective (BWO)

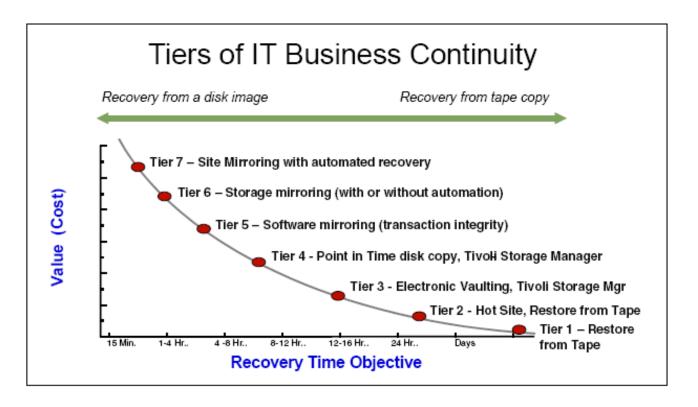
• The *acceptable* amount of time operations can be quiesced to create a copy





Business Continuity Overview (continued)

- Most Business Continuity discussions focus on the *Physical Loss* perspective
 - Power Outages, Fire, Natural Disasters, etc
 - Much money & resources are spent to ensure high RPO and short RTO if such a physical loss should occur

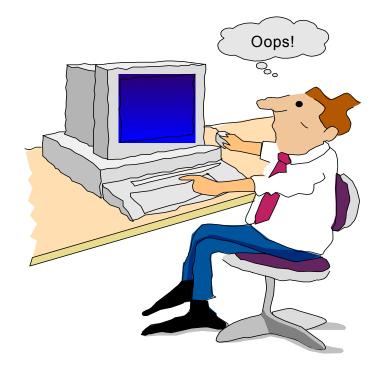






Business Continuity Overview (continued)

- Responsiveness to Logical data loss is sometimes overshadowed
 - Data corruption, User errors, Application errors, Localized data loss, etc
 - What are your RTO / RPO / BWO for these types of data loss?
 - Data-mirroring solutions do not help
 - ! The corruption is just instantly copied from the local site to the remote site
 - Not a question of *if* it will happen, but rather when it will happen



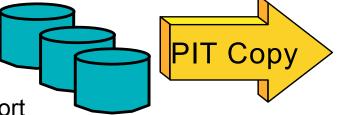


Introduction to Continuous Data Protection

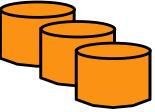
Traditional Point-in-Time Backup

- Taken at specific time or data points
- Only captures data at the point of the backup
- Low RPO
- RTO varies
 - Disk Short
 - Tape Long
- BWO varies
 - Point-in-Time Copy Short
 - Standard I/O Long
- Wouldn't it be nice to be able to recover to a point right before the data was corrupted?









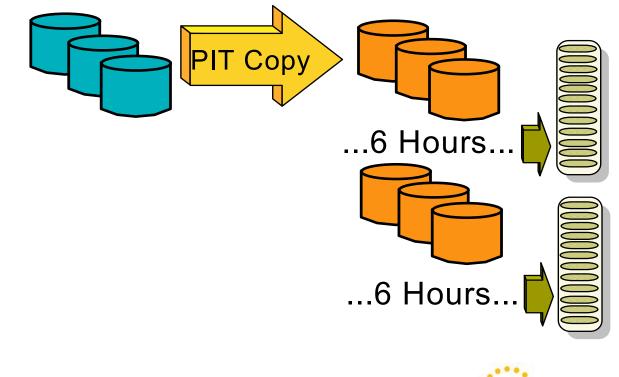
...6 Hours...



Introduction to Continuous Data Protection (continued)

Continuous Data Protection (CDP):

- ★ Continuously captures all changes
 - Journaling combined with Point-in-Time copies
- ★ Eliminates backup window
 - Short/Transparent BWO
- ★ High RPO
- ★ Generally short RTO
 - Long from tape





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Types of Continuous Data Protection

Block Based

- Capture done at storage level
- Time-ordered capture of every block write
- Capturing process does not 'understand' the data
 - Post processing may be required for a data consistent recovery
- True CDP

Application Based

- Specific application journals every update
- Recovery is tightly integrated with the application
- Enables data consistent recovery
- True CDP

File Based

- Runs on application host (Linux, AIX, Windows, etc)
- Backup created when file is written to disk
- Policies can be based on needs of various file types
- Near CDP



zCDP for DB2

- <u>Application based</u> CDP for DB2 on System z
 - Joint solution between DFSMS and DB2
- Solution based on Point-in-Time (PIT) backups combined with DB2 logging
 - ★ Eliminates need for DB2 Log Suspend
 - Only Object-level creates, extends, renames and deletes are suspended
 - Hundreds of volumes backed up in a matter of minutes
 - ★ Managed tape copies created from PIT copies
 - ★ Recovery at the System or Tablespace level
- Base Support: DB2 V8, z/OS V1R5 (2003)
- Enhanced Support: DB2 9, z/OS V1R8, V1R11, V1R12, & V1R13



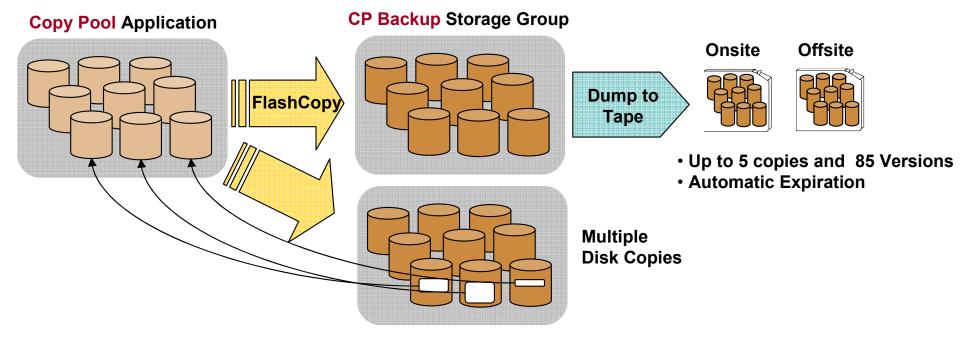


zCDP for DB2



HSM function that manages Point-in-Time copies

• Combined with DB2 BACKUP SYSTEM, provides non-disruptive backup and recovery to any point in time for DB2 databases and subsystems (SAP)



*****Recovery at all levels from either disk or tape!

- Entire copy pool, individual volumes and ...
- Individual data sets



SHARE Technology - Contestions - Results

SMS Constructs

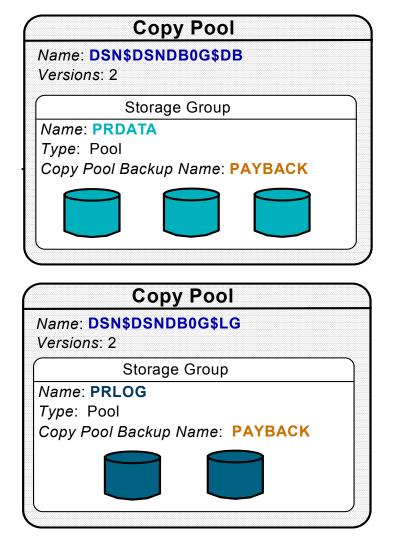
"Copy pool" SMS construct

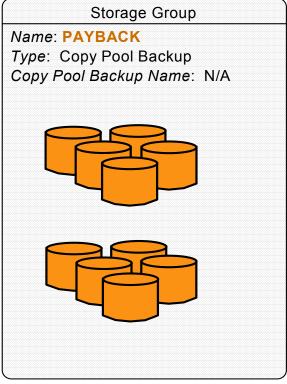
- Defines which storage groups should be processed collectively for point-intime functions
- "Copy pool backup" storage group type
 - Defines which volumes DFSMShsm may select as target volumes for point-in-time backup versions
- Defined via ISMF and Naviquest





SMS Enhancements (continued)





DB2 copy pool names: DSN\$location_name\$DB DSN\$location_name\$LG



S H A R E Technology - Connections - Results

SMS Enhancements (continued)

"copy pool" SMS construct

- Contains from 1 to 256 storage groups
 - Backup processed at storage group level because it is a track-level function
 - Requires data base data to be separated by storage groups
- Specifies the number of *disk* backup copies to maintain
 - 0 85 copies
 - DB2 limit is 50
 - Minimum of 2 copies is recommended
 - Creation of new copy overwrites existing copy
- ★ Storage group volumes dynamically retrieved at time of each backup
 - Ensures that every volume is included in the backup copy



SMS Enhancements (continued)



- "copy pool backup" storage group type
 - Defines candidate target volumes for DFSMShsm fast replication
 - Cannot be assigned for allocation by SMS ACS selection routines
 - ★ Protects data from being overwritten
 - For each source volume to be copied in a storage group:
 - There must exist enough eligible target volumes in copy pool backup storage group to satisfy the needs of the # of specified backup versions



SHARE Technology - Contections - Results

DFSMShsm Support

- DFSMShsm invokes volume-level fast replication to create backup versions for sets of volumes
 - DFSMShsm manages disk and tape copies
- Generic term of "Fast Replication"
 - Multiple Point-in-Time copy functions supported by DFSMSdss
 - FlashCopy
 - SnapShot

Any vendor disk that supports these interfaces





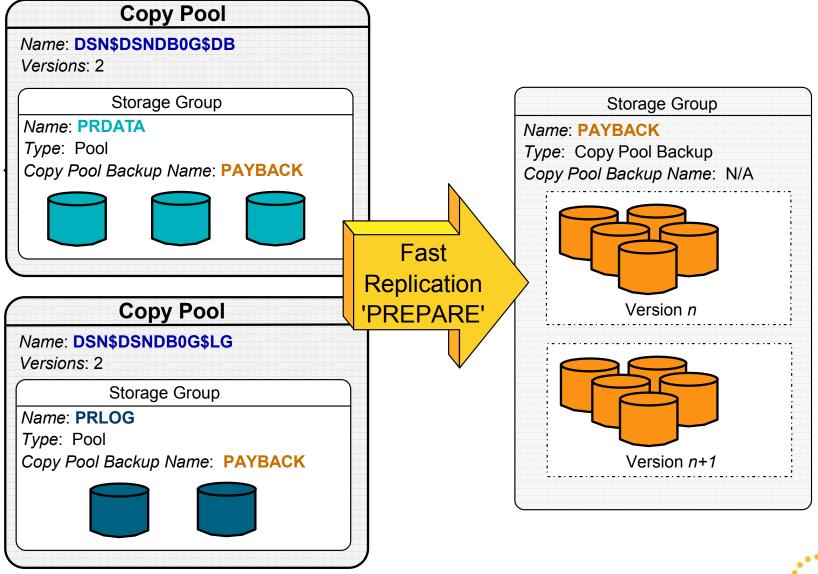
Preparing for Fast Replication Backup

- PREPARE option provided to <u>validate environment</u> and <u>reduce elapsed</u> time of actual backup window
- When PREPARE is specified on FRBACKUP command:
 - For each version > 0, DFSMShsm pre-assigns a target volume to each source volume in every storage group assigned to the copy pool
 - Pairings are maintained in DFSMShsm control data sets
- The PREPARE function *should* be performed whenever there is a change in the environment, such as:
 - Volumes are added to a storage group
 - # of backup versions changes
 - Storage groups are added to the copy pool
- If PREPARE is not performed, target volume selection occurs during backup window





Preparing for Fast Replication Backup (continued)







Creating a Fast Replication Backup

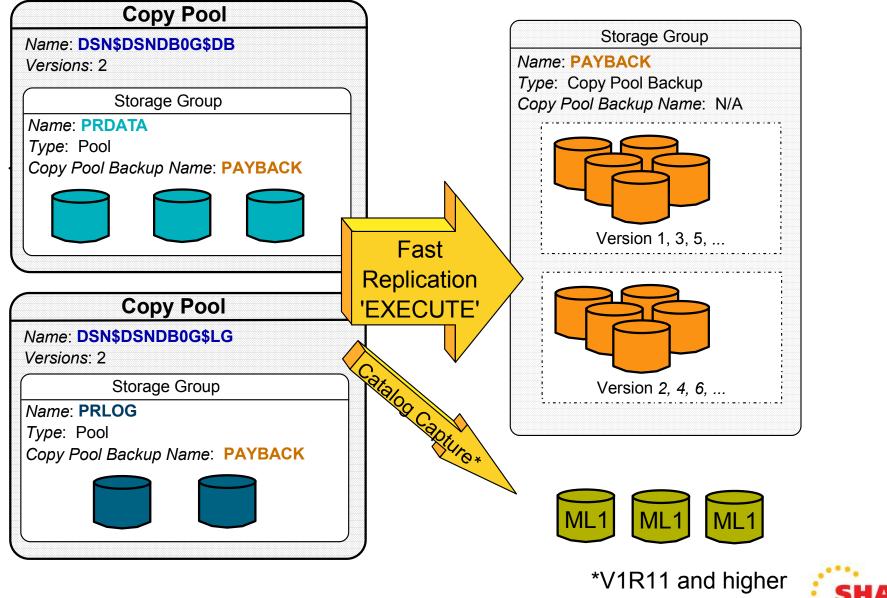
- DFSMShsm command FRBACKUP
 - FRBACKUP COPYPOOL(cpname) EXECUTE TOKEN(token)
- DB2 issues command programmatically via their Utility
- Command is considered successful only after a fast replication relationship has been established for <u>every</u> source volume
- If one or more volumes fail:
 - Version is marked as a failure
 - Failed version will be the target of the next FRBACKUP command
 ★Version number does NOT increment until successfully created





Creating a Fast Replication Backup

(continued)





Creating a Fast Replication Backup

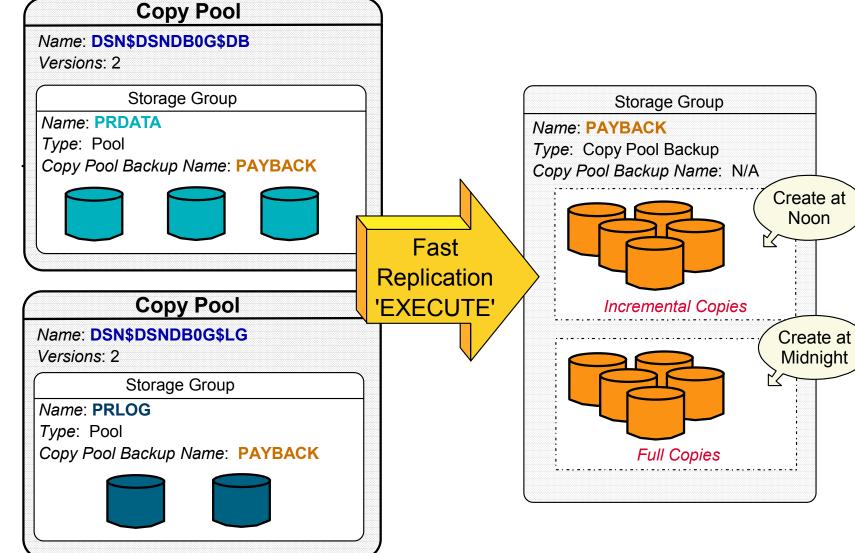


- Versions = 1 to 85
 - Default is to use FlashCopy Full for each source volume
 - FCINCREMENTAL keyword used to specify a volume set for which FlashCopy Incremental should always be used
 - FCINCREMENTAL only needs to be specified once
 - Each time that volume set is Flashed, an Incremental copy is made
- Versions = 0
 - NOCOPY option is used to create FlashCopy
 - Source to Target relationship is Withdrawn after the target volume is dumped to tape
 - ★ Target volumes are available to be used by different copy pools





Creating a Fast Replication Backup





Query



• QUERY COPYPOOL indicates background copy percent complete

ARC18201	THE FOLI	LOWING VO	OLUMES IN COL	PY POOL CP1	, VERSION	003,
HAVE AN ACTIVE FLASHCOPY BACKGROUND COPY						
ARC18201	(CONT.)	SGNAME	FR-PRIMARY	FR-BACKUP	PCT-COMP	
ARC1820I	(CONT.)	SGRP1	SRC01B	TGT01B	70	
ARC1820I	(CONT.)	SGRP1	SRC02B	TGT02B	80	





- **Fast Replication Backup Tape Support**
- DFSMShsm manages copying FlashCopy target volumes to tape
 - By command, immediately after FlashCopy version is created
 - By command, some time after the FlashCopy version is created
 - During Automatic Dump window

Tape copies are 'Dump' copies

- DFSMSdss Full-volume physical dump
- Existing DFSMShsm Dump function used to create copies
 - Dump Classes used to define policies
- Maintain up to 85 versions
 - Up to 5 copies per version





- Dump Class Example
 - Name: ONSITE Copy that is kept onsite Frequency: 0 – Create as often as necessary Retain: 10 Days – Keep for ten days Stacking: 10 – Max volumes on single tape (Higher level of recovery parallelism) Encryption: Yes – Use software encryption on tape
 - Name: OFFSITE Copy that is taken offsite Frequency: 7 – Create once a week Retain: NOLIMIT – Let the copy roll-off Stacking: 255 – 255 volumes on single tape Encryption: Yes
- Recommendation Define new dump classes for fast replication





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Example

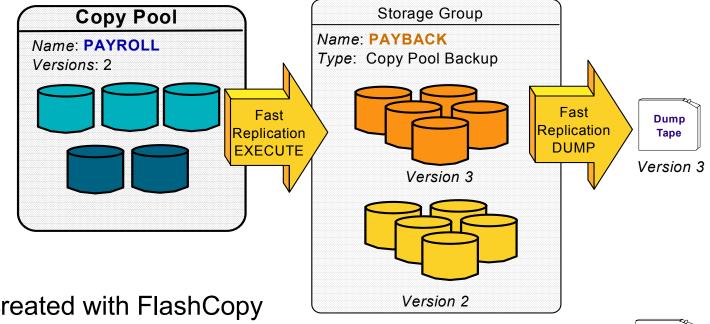
Step 1: Version 1 is created with FlashCopy to disk followed by Dumping target volumes to tape.

Copy Pool Storage Group Name: PAYBACK Name: PAYROLL Type: Copy Pool Backup Versions: 2 Fast Dump Fast Replication Tape Replication DUMP EXECUTE Version 1 Version 1 Unused Storage Group Copy Pool Name: PAYBACK Name: PAYROLL Type: Copy Pool Backup Versions: 2 Fast Replication EXECUTE Version 2 Dump Tape Version 1 Version 1

Step 2: Version 2 is created with FlashCopy to disk only.



Example (continued)



Step 3:

Version 3 is created with FlashCopy to disk followed by Dumping target volumes to tape.

Notice that Version 1 disk copy was overlaid to create Version 3

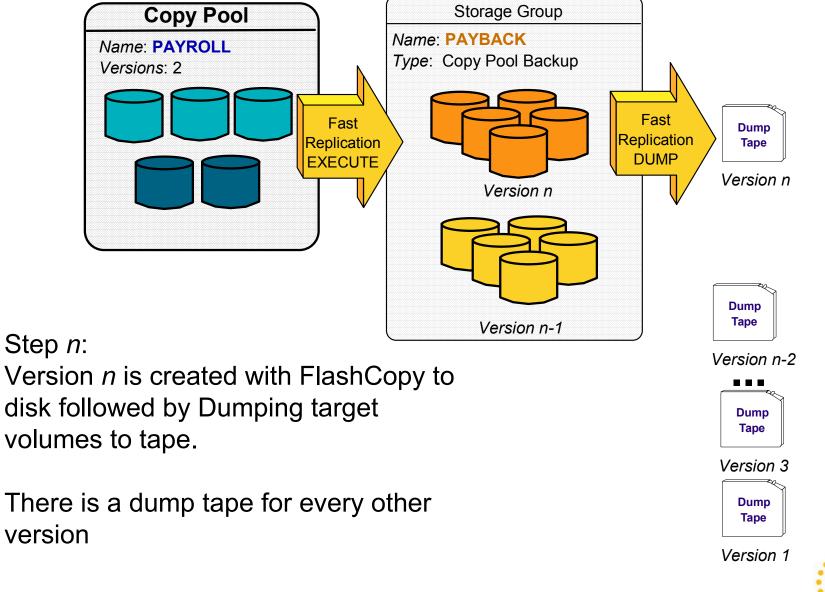


Dump Tape

Version 1



Example (continued)







FRBACKUP COPYPOOL(cpname) DUMP

- After FlashCopy relationships are <u>successfully</u> established for every source volume, DFSMShsm dumps the *target* volumes to tape
- Target volumes are DFSMSdss dump conditioned volumes
 - When dump conditioned volumes are dumped to tape, the tape will look as if the source volume was dumped directly
 - DFSMShsm records will show the dump tape as a dump of the source volume, not the target volume
 - Recovery is done directly back to the source!

Incremental FlashCopy

- Reduces Read I/Os against production volumes for Dump processing
- For performance reasons, some customers wait until background copy is complete before beginning dump copy of target volumes





FRBACKUP COPYPOOL(cpname) DUMPONLY

- Does NOT establish FlashCopy relationships, ONLY creates a dump copy of an <u>existing</u> disk copy
- Use this command to
 - ★Create the dump copy at a time other than when the FlashCopy is created
 - ★Create additional dump copies
 - ★ Retry failed dump copies only dumps previous failures
- Default is to dump Generation 0 Most recent disk copy
 - GENERATION(gennum)
 - VERSION(*vernum*)
 - DATE(date)
 - TOKEN(token)
- The disk copy must be Valid to issue DUMPONLY against it



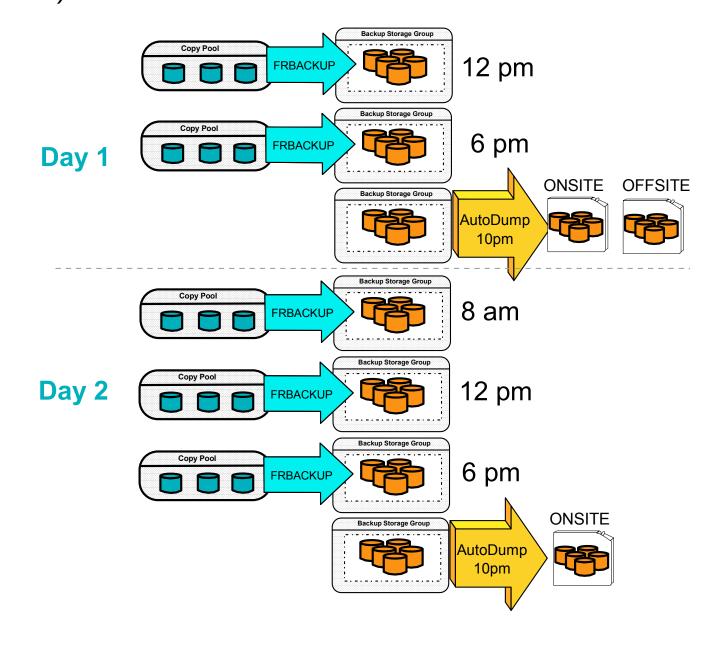


Fast Replication Backup Tape Support (continued)

- Copy Pool definition determines if target volumes are dumped during Automatic Dump window
 - Dumps Generation(0) only
 - If Generation(0) disk copy is failed, then copy pool is not dumped
 - Copy Pool volumes are processed before storage groups and nonSMS volumes
 - System Affinity can be specified in copy pool definition
 - Dump workload is spread across systems
- Using Automatic Dump is a good way to create a single daily tape copy when multiple FlashCopies are created throughout the day











Data Integrity

- Once a dump copy has started, DFSMShsm prevents the target volumes from being overlaid with a new FlashCopy until all volumes have been dumped successfully
 - Overlaying the target volumes in the middle of the dump creates a data integrity exposure.
 - Once the target volumes have been overlaid, it is impossible to complete a dump copy!

DFSMShsm prevents these!

• Withdrawing a background copy while creating a dump copy creates a data integrity exposure

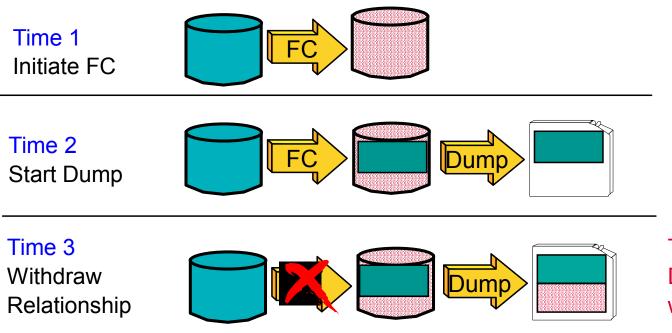
DFSMShsm prevents this!



Data Integrity (continued)



Scenario: Relationship is Withdrawn



Tape is corrupt. Data copied after the withdraw is residual.

DFSMShsm prevents this!

(When Withdraw done with DFSMShsm)





Recovering a Fast Replication Backup

- Target volumes cannot be used as a source volume
 - Dump conditioned volumes
- DFSMShsm FRRECOV command must be used to recover target *disk* volume back to the source
 - Data is immediately available after the FlashCopy initialization completes (High RTO!)
 - ★ DB2 fast log apply can begin within minutes
- FRRECOV COPYPOOL(cpname) VERIFY(Y)
 - Recovers all of the volumes from the named copy pool

★ FRRECOV DSNAME(dsname1, dsname2, ...) REPLACE

- Recovers one or more data sets
- Prefers disk recovery if both disk and tape are available



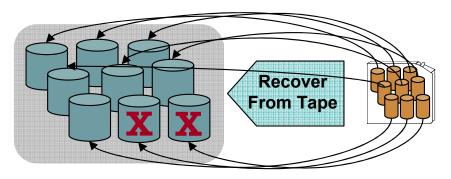


Recovering a Fast Replication Backup

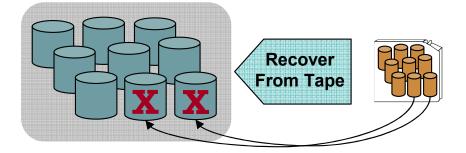
(continued)

Retry Logic

• If there is an error recovering one or more volumes...



- Correct the problem
- Reissue the FRRECOV command
- ★ DFSMShsm will only process those volumes that previously failed
- ★ LIST COPYPOOL output indicates if a particular version recovery is resumable



RESUME(NO) option to retry all volumes (YES) is the default



LIST Command



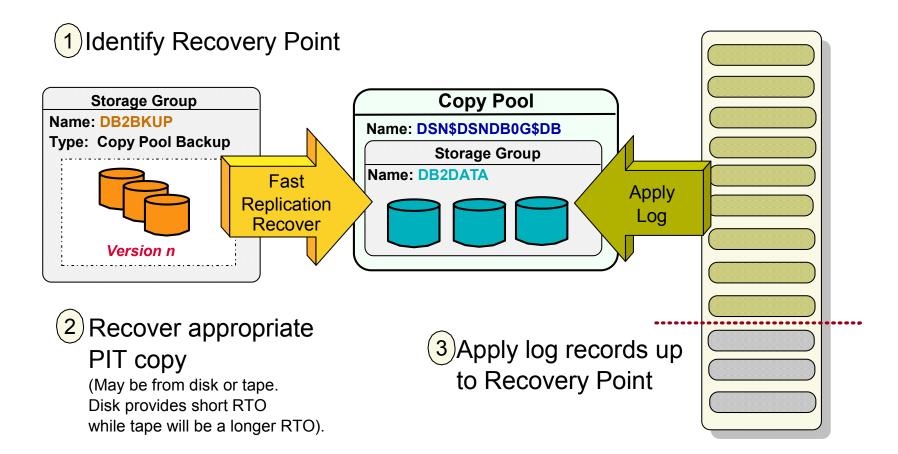
- Returns version information for each copy pool that has at least one attempted backup
- LIST COPYPOOL(cpname)
 - Detailed version information for each backup version for the specified copy pool name
 - Options: <u>FRVOLS</u>, NOVOLS, DUMPVOLS, ALLVOLS
- LIST COPYPOOLBACKUPSTORAGEGROUP(cpbsgname)
 - Information regarding which volumes in copy pool backup storage group that DFSMShsm is using
- SELECT based on state of the version
 - FASTREPLICATIONSTATE
 - Recoverable, NonRecoverable, Failed, None
 - DUMPSTATE
 - AllComplete, RequiredComplete, Partial, None





DB2 RESTORE SYSTEM







Summary



- IBM solution for Continuous Data Protection
 - Cross-Product synergy
- Future enhancements will exploit new advanced features of Disk
- References
 - 'Casebook: DB2 backup, recovery and cloning for SAP environments'

https://www.sdn.sap.com/irj/scn/go/portal/prtroot/docs/library/uuid/e0b13d2b-0a89-2b10-918b-bc7aff0a1905

- <u>z/OS Host Topics</u> Issue 16, February 2007
- Redbook: DFSMShsm Fast Replication Technical Guide

